# A LOW COST IMPLEMENTATION OF ADVANCED ENCRYPTION STANDARD ALGORITHM USING 8085A MICROPROCESSOR 

Salim M. Wadi ${ }^{1 *}$, Nasharuddin Zainal ${ }^{1 \dagger}$,<br>${ }^{1}$ Department of Electrical, Electronic and System Engineering , UKM University, Malaysia


#### Abstract

The high security communication systems became an urgent need in recent years for both governments and peoples desiring protection from signal interception. The implementation of advanced encryption standard algorithm is important requirement where many researches proposed different items to this purpose. Some papers used microcontrollers as CPU item to implement AES algorithm. A simply item proposed in this paper to speedy, low cost implementation of Rijndael Advanced Encryption Standard (AES) cryptographic algorithm which is 8085A microprocessor. The results prove the implementation is effective through competitive cost, low hardware requirements, and reasonable speeds compared with other implementation methods.


Keywords: AES, 8085 processor, ALP, microcontroller, Rijndael.

## 1. INTRODUCTION

Transmission of important electronic financial transactions and digital touch applications must be fast and very secure to achieve the requirements for security, integrity, and non reproduction of exchanged information. Cryptography provides a method for securing and authenticating the transmission of information over insecure channels. For these reasons, large number of research were done to developing a high performance encryption system (Daemen and Rijmen 2002; Tran et al 2008) .

Vincent Rijmen and Joan Daeman were innovated new encryption system which referred to as advanced encryption standard (AES) in 2001(Daemen and Rijmen 2002). The Advanced Encryption Standard (AES) was successor alternative algorithm to the Data Encryption Standard (DES) which suffers from theoretical weaknesses in the encryption as well as successful brute force attacks carried out against the algorithm(Orlando et al 2008). The implementations of AES are carryout into two ways, one by FPGA and other by microcontroller.

[^0]Orlando, J.et al. (Orlando et al 2008) used a bit-serial approach to built their architecture utilizing FPGA with focusing on low cost resulted in a design well-suited for SoC implementations. Proposed implementation give good results in term throughput/slice ratio and cost can be reduced while maintaining a suitable operating speed with minimization of redundancies.

Hyubgun Lee et al. (Hyubgun et al 2009 ) were presented the sensor network with high security to analyze the communication efficiency through performance evalution of AES ciphering system depending on data length, and cost of operation per hop according to the network scale. The authors concludes that if the scale of the sensor network increased, this lead to doubled the delay as well as increasing in the energy disbursed.

Kai Schramm et al. (Kai and Christ 2004) were design an particularistic lab project in IT security which combines topics of computer architecture, cryptography and software engineering. The professed that undergraduate Students of EE/CS are used the proposed lab to worthily carryout the Advanced Encryption Standard (AES) on a smart card using Atmel ATMega163 Reduced Instruction Set Computer (RISC) microcontroller in assembly.

Gielata et al. (Gielata et al 2008) were proposed special architecture using FPGA to get speedy and flexibility implementation to 128 -ASE algorithm. The hardware performance was been evaluated depending on four parameters as, number of used resources, maximum clock frequency, latency measured in number of clocks, throughput measured in bytes per second.

In this paper simply item proposed to low cost implementation of AES algorithm using 8085 microprocessor as processing tool with very effective manner and reasonable speed.
the paper is organized as, in section2 clarify short for AES algorithm. Explain to implementation method show in section3. Results and comparison in section 4 and conclusion in section 5.

## 2- AES ALGORITHM

AES is based on Rijndael algorithm which is a symmetric block cipher that processes fixed data of 128 -bit blocks. It supports key sizes of 128,192 and 256 bits and consists of 10,12 or 14 iteration rounds, respectively. In this paper, we will present the 128 -bit version of AES with 10 rounds. Each round mixes the data with a round key, which is generated from the encryption key. The AES encryption structure is shown in figure 1 (Orlando et al 2008; Schmidt et al 2009). The cipher maintains an internal, $4 \times 4$ matrix of bytes referred to as State, on which the operations are performed. Initially, State is filled with the input data block and XOR-ed with the encryption key. Regular rounds consist of operations called SubBytes, ShiftRows, MixColumns and AddRoundKey. The last round bypasses MixColumns transformation. SubBytes transformation uses 16 identical 256-byte substitution table called S-box as shown in table1(Majithia and Dinesh 2010).

SubBytes can be implemented either by computing the substitution or using look-up-table (LUT). ShiftRows is a cyclic left shift of the second, third and fourth row of State by one, two, and three bytes, respectively. MixColumns performs a modular polynomial multiplication on each column.


Figure 1: AES Encryption Structure.
During each round, AddRoundKey performs XOR with State and the round key. Round key generation (key expansion) includes S-box substitutions, word rotations, and XOR operations performed on the encryption key. Depending on the security level required for the application, AES uses different key length (Muhammad and Syed 2009; Ming et al 2007).

Table 1: S-Box based on Galois Field GF ( $\mathbf{2}^{\mathbf{8}}$ )

|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | a | b | c | d | e | f |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 63 | 7 c | 77 | 7b | f2 | 6b | 6 f | c5 | 30 | 01 | 67 | 2b | fe | d7 | ab | 76 |
| 1 | ca | 82 | c9 | 7 d | fa | 59 | 47 | f0 | ad | d4 | a2 | af | 9 c | a4 | 72 | c0 |
| 2 | b7 | fd | 93 | 26 | 36 | 3f | f7 | cc | 34 | a5 | e5 | f1 | 71 | d8 | 31 | 15 |
| 3 | 04 | c7 | 23 | c3 | 18 | 96 | 05 | 9a | 07 | 12 | 80 | e2 | eb | 27 | b2 | 75 |
| 4 | 09 | 83 | 2c | 1a | 1b | 6 e | 5a | a0 | 52 | 3b | d6 | b3 | 29 | e3 | 2f | 84 |
| 5 | 53 | d1 | 00 | ed | 20 | fc | b1 | 5 b | 6a | cb | be | 39 | 4a | 4 c | 58 | cf |
| 6 | d0 | ef | aa | fb | 43 | 4d | 33 | 85 | 45 | f9 | 02 | 7 f | 50 | 3c | 9 f | a8 |
| 7 | 51 | a3 | 40 | 8 f | 92 | 9d | 38 | f5 | bc | b6 | da | 21 | 10 | ff | f3 | d2 |
| 8 | cd | Oc | 13 | ec | 5 f | 97 | 44 | 17 | c4 | a7 | 7 e | 3d | 64 | 5d | 19 | 73 |
| 9 | 60 | 81 | 4 f | dc | 22 | 2a | 90 | 88 | 46 | ee | b8 | 14 | de | 5e | 0b | db |
| a | e0 | 32 | 3 a | 0a | 49 | 06 | 24 | 5 c | c2 | d3 | ac | 62 | 91 | 95 | e4 | 79 |
| b | e7 | c8 | 37 | 6d | 8d | d5 | 4 e | a9 | 6 c | 56 | f4 | ea | 65 | 7a | ae | 08 |
| c | ba | 78 | 25 | 2e | 1 c | a6 | b4 | c6 | e8 | dd | 74 | 1f | 4b | bd | 8b | 8 a |
| d | 70 | 3 e | b5 | 66 | 48 | 03 | f6 | 0 0 | 61 | 35 | 57 | b9 | 86 | c 1 | 1d | 9 e |
| e | el | 88 | 98 | 11 | 69 | d9 | 8 e | 94 | 9 b | 1e | 87 | e9 | ce | 55 | 28 | df |
| f | 8 c | a 1 | 89 | Od | bf | e6 | 42 | 68 | 41 | 99 | 2 d | Of | b0 | 54 | bb | 16 |

## 3- PROPOSED IMPEMENTATION:

The best encryption algorithm that is easy to implement in software and in hardware (Orlando et al 2008). Inventors of AES algorithm designed it with an idea in mind of ability for its efficient execution using different platform such as CPU, ASICs, and FPGA. Microprocessor is cheap and easy programmed and has high efficiency. For these reasons, we choose the microprocessor to implement the AES algorithm. A 8085 microprocessor is used to implement AES algorithm in this paper because it compatible with proposed requirements (low cost, good speed, and efficiency). In this implementation, SubBytes transform implemented using lookup table (LUT). MixColumns transform implemented using XOR operations with mod to (11B) because the state in hexadecimal numbers. ShiftRows and AddRoundKey transforms are easily executed using 8085 microprocessor.

A block diagram to proposed implementation shown in figure 2, which consist from:

- 8085A microprocessor: this processor is from Intel family, 8 bit processor, operating frequency is 5 MHz .
- EEPROM: about 0.978 Kbyte used to save the operating programs and S Box and inverse S Box.
- RAM: about 0.76 Kbyte
- Keyboard and 7-sigment display to input the plaintext and key block as inputs and display the cipher text as outputs.
- I/O ports.


Figure 2: Block diagram of device.

## 4- RESULT:

Performance of the proposed implementation was evaluated according to the below parameters: Simplicity: the proposed implementation is very simple as hardware see fig. 2 as well as the programming requirements is very easy where assembly language used to write the essential programs to employment the microprocessor.
Cost: As known to all, the price of 8085 microprocessor IC is very low about 1-1.5 \$ which is the mean part of the proposed circuit and this is true for other parts, therefore we think that the cost of the complete circuit is very low compared with other implementation using FPGA or microcontrollers.
Speedy: Table 2 shown the implementation results to the circuit. From this table we see that the speed of complete implementation is reasonable and compatible most applications as example to researches purpose.

Table 2: The implementation results

| Implementation of AES | Suggested implement. |
| :--- | :--- |
| Type of processor | $8085 \mathrm{~A} \mu \mathrm{P}$ |
| Frequency | 5 MHz |
| Time of encryption | 17.66 ms |
| No. of CPU Cycle (for one block data) | 4500 |
| EEPROM needed | 0.978 Kbyte (contain the program memory) |
| Volatile memory | 0.76 Kbyte |
| Internal registers | $7(8-$-bit), 1 (16-bit) |

To demonstrate the effectiveness of proposed method, we compare the proposed implementation with other two implementations, which used the microcontroller. We don't find implementation using any type of microprocessor especially 8085 microprocessor. Table 3 shown the comparison.

From table 3, we see that the consumption time in proposed method near of the proposed method in (Kai and Christ 2004) but the memory required is little in proposed method and is cheaper. In compare with implementation that proposed in (Hyubgun et al 2009) the difference in terms of time carryout and memory required is very high. Therefore, we conclude that the proposed method is effective in terms of cost, time consumption, simplicity, and required peripherals such as memory.

Table 3:Comparison of Suggested Implementation Results with Proposed Implementation in [4, 5]

| Implementation of AES | Suggested implement. | Prop. imp. in [4] | Prop. imp. in [5] |
| :--- | :--- | :--- | :--- |
| Type of processor | $8085 \mathrm{~A} \mu \mathrm{P}$ | ATmega644p $\mu \mathrm{C}$ | ATMega163 $\mu \mathrm{C}$ |
| Frequency | 5 MHz | 12 MHz | 3.57 MHz |
| Time of encryption | 17.66 ms | 449 ms | 22.50 ms |
| EEPROM needed | 0.978 Kbyte (contain the <br> program memory) | 8 Kbyte | 32 byte |
| Volatile memory | 0.76 Kbyte | Not found | 2.234 Kbyte |
| Internal registers | $7(8$-bit), 1 (16-bit) | Not found | 32 (8-bit) |

## 5-CONCLUSION:

In this paper we proposed the novel implementation to AES algorithm using 8085 microprocessor. The 8085 microprocessor is low cost and effectiveness with reasonable speed. the performance of proposed implementation is evaluated in terms of cost and time required to carry out the AES algorithm. As a result the time and cost of the proposed implementation is good in compare with other implementation.

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## Appendix A Example of AES Encryption

Table1: key expansion example

| Key words | Auxiliary function |
| :---: | :---: |
| W0=0f 1571 c 9 <br> W1=47 d9 e8 59 <br> W2=0c b7 ad d6 <br> W3=af 7f 6798 | Rot word (W3)=7f 6798 af $=x 1$ <br> Sub word (x1) =d2 $8546 \quad 79=y 1$ <br> $R$ con (1) $=01000000$ <br> $\mathrm{Y} 1 \oplus \mathrm{R} \operatorname{con}(1)=\mathrm{d} 3854679=\mathrm{z} 1$ |
| $\begin{aligned} & \mathrm{W} 4=\mathrm{W} 0 \oplus \mathrm{Z} 1=\begin{array}{lllll} \text { dc } & 90 & 37 & \mathrm{~b} 0 \\ \mathrm{~W} 5=\mathrm{W} 4 \oplus \mathrm{~W} 1=9 \mathrm{~b} & 49 & \text { df } & \mathrm{e} 9 \end{array} \\ & \mathrm{~W} 6=\mathrm{W} 5 \oplus \mathrm{~W} 2=\begin{array}{lllll} 97 & \text { fe } & 72 & 3 \mathrm{f} \end{array} \\ & \mathrm{~W} 7=\mathrm{W} 6 \oplus \mathrm{~W} 3=38 \\ & 31 \end{aligned} 15 \mathrm{a} 7 .$ | Rot word (W7)=81 15 a7 $38=x 2$ <br> Sub word $(x 2)=0 c \quad 595 c \quad 07=y 2$ <br> $R$ con $(2)=02000000$ <br> $\mathrm{Y} 2 \oplus \mathrm{R} \operatorname{con}(2)=\mathrm{d} 38546 \quad 79=\mathrm{z} 2$ |
|  | Rot word (W11) $=\mathrm{ff}$ d3 c6 e6 $=\mathrm{x} 3$ <br> Sub word $(x 3)=1666$ b4 8e $=y 3$ <br> $R$ con (3) $=04000000$ <br> $\mathrm{Y} 3 \oplus \mathrm{R} \operatorname{con}(3)=1266 \mathrm{~b} 48 \mathrm{e}=\mathrm{z} 3$ |
| $\begin{aligned} & \mathrm{W} 12=\mathrm{W} 8 \oplus \mathrm{Z} 3=\mathrm{c} 0 \text { af df } 39 \\ & \mathrm{~W} 13=\mathrm{W} 12 \oplus \mathrm{~W} 9=89 \text { 2f } 6 \mathrm{~b} 67 \\ & \mathrm{~W} 14=\mathrm{W} 13 \oplus \mathrm{~W} 10=5751 \text { ad } 06 \\ & \mathrm{~W} 15=\mathrm{W} 14 \oplus \mathrm{~W} 11=\mathrm{b} 1 \text { ae } 7 \mathrm{e} c 0 \end{aligned}$ | Rot word (W15) $=$ ae $7 \mathrm{e} \mathrm{c} 0 \quad \mathrm{~b} 1=\mathrm{x} 4$ <br> Sub word $(x 4)=e 4$ f3 ba c8 $=y 4$ <br> $R$ con $(4)=08000000$ <br> $\mathrm{Y} 4 \oplus \operatorname{Rcon}(4)=\mathrm{ec} \mathrm{f} 3 \mathrm{ba} \mathrm{c} 8=\mathrm{z} 4$ |
|  | Rot word (W19) $=8 \mathrm{c}$ dd $5043=x 5$ <br> Sub word (x5) $=64$ c1 $53 \quad 1 \mathrm{a}=\mathrm{y} 5$ <br> $R$ con (5) $=10000000$ <br> $\mathrm{Y} 5 \oplus \mathrm{R} \operatorname{con}(5)=74 \mathrm{c} 153 \quad 1 \mathrm{a}=\mathrm{z} 5$ |
| $\begin{aligned} & \mathrm{W} 20=\mathrm{W} 16 \oplus \mathrm{Z} 5=589 \mathrm{~d} 36 \mathrm{eb} \\ & \mathrm{~W} 21=\mathrm{W} 20 \oplus \mathrm{~W} 17=\text { fd ee } 387 \mathrm{~d} \\ & \mathrm{~W} 22=\mathrm{W} 21 \oplus \mathrm{~W} 18=0 \mathrm{fc} 9 \mathrm{~b} \text { ed } \\ & \mathrm{W} 23=\mathrm{W} 22 \oplus \mathrm{~W} 19=4 \mathrm{c} 4046 \mathrm{bd} \end{aligned}$ | Rot word $(W 23)=4046$ bd $4 \mathrm{c}=\mathrm{x} 6$ <br> Sub word (x6) =09 5a 7a $29=y 6$ <br> $R$ con (6) $=20000000$ <br> $\mathrm{Y} 6 \oplus \mathrm{R} \operatorname{con}(6)=295 \mathrm{a} 7 \mathrm{a} 29=\mathrm{z} 6$ |
|  | $\begin{array}{llll} \operatorname{Rot} \operatorname{word}(\mathrm{W} 27) & =a 5 & \mathrm{a} 9 & \text { ef } \mathrm{cf}=\mathrm{x} 7 \\ \text { Sub word }(\mathrm{x} 7) & =06 & \mathrm{~d} 3 & \mathrm{df} \\ 8 \mathrm{a}=\mathrm{a} 7 \\ \mathrm{R} \operatorname{con}(7) & =40 & 00 & 00 \\ \mathrm{Y} 7 \oplus \operatorname{Rcon}(7) & =46 & \mathrm{~d} 3 & \text { df } 8 \mathrm{a}=\mathrm{z} 7 \end{array}$ |
|  |  |
|  |  |
| $\begin{aligned} & \mathrm{W} 36=\mathrm{W} 32^{\oplus} \mathrm{Z} 9=\mathrm{fd} \\ & \mathrm{~W} 9 \mathrm{~d} 42 \mathrm{cb} \\ & \mathrm{~W} 37=\mathrm{W} 36^{\oplus} \mathrm{W} 33=0 \mathrm{e} \\ & \hline \end{aligned} 6 \mathrm{e} 0 \quad 1 \mathrm{c}$ |  |

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W40= W36 © Z10 = b4 8e f3 52
W41= W40 }\oplus\textrm{W}37=\mathrm{ ba 98 13 4e
W42= W41 \oplus W38 = 7f 4d 59 20
W43=W42 }\oplus\textrm{W}39=8626187
```

Table 2: AES encryption example

| Start of round | After sub Byte | After Shift Ro. | After Mix col. | Round Key |
| :---: | :---: | :---: | :---: | :---: |
| 0189 fe 76 <br> 23 ab dc 54 <br> 45 cd ba 32 <br> 67 ef $98 \quad 10$ |  |  |  | Of 47 $0 c$ af <br> 15 d9 b7 7 f <br> 71 e8 ad 67 <br> c9 59 d6 98 |
| $\begin{array}{cccc} 0 \mathrm{e} & \text { ce } & \text { f2 } & \text { d9 } \\ 36 & 72 & 6 \mathrm{~b} & 2 \mathrm{~b} \\ 34 & 25 & 17 & 55 \\ \text { ae } & \text { b6 } & 4 \mathrm{c} & 88 \end{array}$ | $\begin{array}{cccc} \mathrm{ab} & 8 \mathrm{~b} & 89 & 35 \\ 05 & 40 & 7 \mathrm{f} & \mathrm{f} 1 \\ 18 & 3 \mathrm{f} & \mathrm{fo} & \mathrm{fc} \\ \mathrm{e} 4 & 4 \mathrm{e} & 2 \mathrm{f} & \mathrm{c} 4 \end{array}$ | ab 8 8b 8935 <br> 40 7f f1 05 <br> f0 fc 18 3f <br> c4 $\mathrm{e} 4 \quad 4 \mathrm{e} \quad 2 \mathrm{f}$ | b9 945775 <br> c4 8e 1651 <br> 47 20 9a 3f <br> c5 d6 f5 3b | $\begin{array}{cccc} \text { dc } & 9 \mathrm{~b} & 97 & 38 \\ 90 & 49 & \text { fe } & 81 \\ 37 & \text { df } & 72 & 15 \\ \text { b0 } & \text { e9 } & 3 \mathrm{f} & \text { a7 } \end{array}$ |
| $\begin{array}{llll} 65 & 0 f & \text { c0 } & 4 d \\ 74 & \mathrm{c} 7 & \mathrm{e} 8 & \mathrm{~d} 0 \\ 70 & \text { ff } & \text { e8 } & 2 \mathrm{a} \\ 75 & 3 \mathrm{f} & \mathrm{ca} & 9 \mathrm{c} \end{array}$ | $4 d$ 76 ba $e 3$ <br> 92 $c 6$ $9 b$ 70 <br> 51 16 $9 b$ $e 5$ <br> $9 d$ 75 74 de | $\begin{array}{cccc} 4 d & 76 & \text { ba } & \text { e3 } \\ \text { c6 } & 9 b & 70 & 92 \\ 9 b & \text { e5 } & 51 & 16 \\ \text { de } & 9 d & 75 & 74 \end{array}$ | $\begin{array}{cccc} 8 \mathrm{e} & 22 & \mathrm{db} & 12 \\ \mathrm{~b} 2 & \mathrm{f} 2 & \mathrm{dc} & 92 \\ \mathrm{df} & 80 & \mathrm{f} 7 & \mathrm{c} 1 \\ 2 \mathrm{~d} & \mathrm{c} 5 & 1 \mathrm{e} & 52 \end{array}$ | d2 49 de e6 <br> c9 50 7e ff <br> 6b b4 c6 d3 <br> b7 $5 \mathrm{Ee} 61 \quad$ c6 |
| $\begin{array}{cccc} 5 \mathrm{c} & 6 \mathrm{~b} & 05 & \mathrm{f} 4 \\ 7 \mathrm{~b} & 72 & \mathrm{a} 2 & 6 \mathrm{~d} \\ \mathrm{~b} 4 & 34 & 31 & 12 \\ 9 \mathrm{a} & 9 \mathrm{~b} & 7 \mathrm{f} & 94 \end{array}$ | $\begin{array}{cccc} 4 \mathrm{a} & 7 \mathrm{f} & 6 \mathrm{~b} & \mathrm{bf} \\ 21 & 40 & 3 \mathrm{a} & 3 \mathrm{c} \\ 8 \mathrm{~d} & 81 & \mathrm{c} 7 & \mathrm{c} 9 \\ \mathrm{~b} 8 & 14 & \mathrm{~d} 2 & 22 \end{array}$ | $\begin{array}{llll} 4 \mathrm{a} & 7 \mathrm{f} & 6 \mathrm{~b} & \mathrm{bf} \\ 40 & 3 \mathrm{a} & 3 \mathrm{c} & 21 \\ \mathrm{c} 7 & \mathrm{c} 9 & 8 \mathrm{~d} & 81 \\ 22 & \mathrm{~b} 8 & 14 & \mathrm{~d} 2 \end{array}$ | b1 c1 0b cc ba f3 8b 07 f9 1f 6a c3 1 1d $19 \quad 24$ 5c | $\begin{array}{llll}\text { c0 } & 89 & 57 & \text { b1 }\end{array}$ <br> af 2 f 51 ae <br> df 6 b ad 7 e <br> $3967 \quad 06 \quad c 0$ |
| $\begin{array}{cccc} 71 & 48 & 5 \mathrm{c} & 7 \mathrm{~d} \\ 15 & \mathrm{dc} & \mathrm{da} & \text { a } 9 \\ 26 & 74 & \mathrm{c} 7 & \mathrm{bd} \\ 24 & 7 \mathrm{e} & 22 & 9 \mathrm{c} \end{array}$ | a3 52 4 a ff <br> 59 86 57 d3 <br> f7 92 c6 7 a <br> 36 f3 93 de | $\begin{array}{llll} \text { a3 } & 52 & 4 \mathrm{a} & \text { ff } \\ 86 & 57 & \text { d3 } & 59 \\ \text { c6 } & 7 \mathrm{a} & \text { f7 } & 92 \\ \text { de } & 36 & \text { f3 } & 93 \end{array}$ | d4 11 fe $0 f$ <br> 3b $44 \quad 0673$ <br> cb ab 6237 <br> 19 b7 07 ec | $\begin{array}{cccc} 2 \mathrm{c} & \text { a5 } & \mathrm{f} 2 & 43 \\ 5 \mathrm{c} & 73 & 22 & 8 \mathrm{c} \\ 65 & 0 \mathrm{e} & \text { a3 } & \text { dd } \\ \text { f1 } & 96 & 90 & 50 \end{array}$ |
| $\begin{array}{llll}\text { f8 } & \mathrm{b} 4 & 0 \mathrm{c} & 4 \mathrm{c}\end{array}$ $\begin{array}{lll}67 & 37 & 24\end{array}$ ae a5 c1 ea e8 2197 bc | $\begin{array}{llll} 41 & 8 \mathrm{~d} & \text { fe } & 29 \\ 85 & 9 \mathrm{a} & 36 & 16 \\ \text { e4 } & 06 & 78 & 87 \\ 9 \mathrm{~b} & \text { fd } & 88 & 65 \end{array}$ | $\begin{array}{llll} 41 & 8 \mathrm{~d} & \text { fe } & 29 \\ 9 \mathrm{a} & 36 & 16 & 85 \\ 78 & 87 & \text { e4 } & 06 \\ 65 & 9 \mathrm{~b} & \text { fd } & 88 \end{array}$ | $\begin{array}{llll} 2 \mathrm{a} & 47 & \mathrm{c} 4 & 48 \\ 83 & \mathrm{e} 8 & 18 & \mathrm{ba} \\ 84 & 18 & 27 & 23 \\ \mathrm{eb} & 10 & 0 \mathrm{a} & \mathrm{f} 3 \end{array}$ | 58 fd of 4 c <br> 9 d ee cc 40 <br> $36 \quad 38 \quad 9 b 46$ <br> eb 7d ed bd |
| $\begin{array}{llll} 72 & \mathrm{ba} & \mathrm{cb} & 04 \\ 1 \mathrm{e} & 06 & \mathrm{~d} 4 & \mathrm{fa} \\ \mathrm{~b} 2 & 20 & \mathrm{bc} & 65 \\ 00 & 6 \mathrm{~d} & \mathrm{e} 7 & 4 \mathrm{e} \end{array}$ | 40 f4 1f f2 <br> 72 6f 48 2d <br> 37 b7 654 d <br> 63 3c 94 2f | $\begin{array}{llll} 40 & \mathrm{f} 4 & 1 \mathrm{f} & \mathrm{f} 2 \\ 6 \mathrm{f} & 48 & 2 \mathrm{~d} & 72 \\ 65 & 4 \mathrm{~d} & 37 & \mathrm{~b} 7 \\ 2 \mathrm{f} & 63 & 3 \mathrm{c} & 94 \end{array}$ | $\begin{array}{cccc} 7 \mathrm{~b} & 05 & 42 & 4 \mathrm{a} \\ 1 \mathrm{e} & \mathrm{~d} 0 & 20 & 40 \\ 94 & 83 & 18 & 52 \\ 94 & \mathrm{c} 4 & 43 & \mathrm{fb} \end{array}$ | 718 c 83 cf <br> c7 29 e5 a5 <br> 4c 74 ef a9 <br> c2 bf 52 ef |
| $\begin{array}{llll} 0 \mathrm{a} & 89 & \mathrm{c} 1 & 83 \\ \text { d9 } & \text { f9 } & \text { c5 } & \text { e5 } \\ \text { d8 } & \text { f7 } & \text { f7 } & \text { fb } \\ 56 & 7 \mathrm{~b} & 11 & 14 \end{array}$ | $\begin{array}{cccc} 67 & \text { a7 } & 78 & 97 \\ 35 & 99 & \text { a6 } & \text { d9 } \\ 61 & 68 & 68 & 0 f \\ \text { b1 } & 21 & 82 & \text { fa } \end{array}$ | $\begin{array}{llll} 67 & \text { a7 } & 78 & 97 \\ 99 & \text { a6 } & \text { d9 } & 35 \\ 68 & \text { of } & 61 & 68 \\ \text { fa } & \text { b1 } & 21 & 82 \end{array}$ | $\begin{array}{cccc} \text { ec } & \text { 1e } & \mathrm{c} 0 & 80 \\ 0 \mathrm{c} & 50 & 53 & \mathrm{c} 7 \\ 3 \mathrm{~b} & \mathrm{~d} 7 & 00 & \text { ef } \\ \text { b7 } & 22 & 72 & \mathrm{e} 0 \end{array}$ | $\begin{array}{llll} 37 & \mathrm{bb} & 38 & \mathrm{f} 7 \\ 14 & 3 \mathrm{~d} & \mathrm{~d} 8 & 7 \mathrm{~d} \\ 93 & \mathrm{e} 7 & 08 & \mathrm{a} 1 \\ 48 & \mathrm{f} & \mathrm{a} 5 & 4 \mathrm{a} \end{array}$ |
| db al f8 77 18 6d 8b ba $\begin{array}{llll}\text { a8 } & 30 & 08 & 4 \mathrm{e}\end{array}$ | $\begin{array}{llll} \mathrm{b} 9 & 32 & 41 & \mathrm{f} 5 \\ \mathrm{ad} & 3 \mathrm{c} & 3 \mathrm{~d} & \mathrm{f} 4 \\ \mathrm{c} 2 & 04 & 30 & 2 \mathrm{f} \\ \hline \end{array}$ | $\begin{array}{cccc} \mathrm{b} 9 & 32 & 41 & \mathrm{f} 5 \\ 3 \mathrm{c} & 3 \mathrm{~d} & \mathrm{f} 4 & \mathrm{ad} \\ 30 & 2 \mathrm{f} & \mathrm{c} 2 & 04 \\ \hline \end{array}$ | $\begin{array}{llll}\text { b1 } & 1 \mathrm{a} & 44 & 17\end{array}$ <br> 3d 2 f ec b6 <br> $0 \mathrm{a} \quad 6 \mathrm{~b} \quad 2 \mathrm{f} \quad 42$ | $\begin{array}{llll} 48 & \mathrm{f} & \mathrm{cb} & 3 \mathrm{c} \\ 26 & \mathrm{bb} & \mathrm{c} 3 & \mathrm{be} \\ 45 & \mathrm{a} 2 & \text { aa } & 0 \mathrm{~b} \\ \hline \end{array}$ |


| ff d5 d7 aa | 16030 e ac | ac $16 \quad 030 \mathrm{e}$ | 9f 68 f3 b1 | 20 d7 $72 \quad 38$ |
| :---: | :---: | :---: | :---: | :---: |
| f9 e9 8f 2 b | 99 1e 73 f1 | 99 1e 73 fl | 31303 a c2 | fd 0 e c5 f9 |
| 1 b | af 1818 | $\begin{array}{llll}18 & 15 & 30\end{array}$ | ac 7188 c c4 | 0d 16 d5 6b |
| 4f c9 80549 | 84 dd 97 3b | 97 3b 84 dd | 466548 eb | 42 e0 4a 41 |
| bf bf 8189 | $08 \quad 08 \quad 0 \mathrm{c}$ a7 | a7 $0808080 c$ | 6a 1c 3162 | cb 1c 6e 56 |
| cc 3e ff 3b | $\begin{array}{lllll}4 b & \mathrm{~b} 2 & 16 & \mathrm{e} 2\end{array}$ | 4b b2 16 e2 |  | b4 8e f3 52 |
| a1 6759 af | 3285 cb 79 | 85 cb 7932 |  | ba $98 \quad 13$ 4e |
| $\begin{array}{llll}04 & 85 & 02\end{array}$ | f2 $97 \begin{array}{llll}77 & \mathrm{ac}\end{array}$ | 77 ac f2 97 |  | 7f 4d $59 \quad 20$ |
| a1 00 5f 34 | $32 \quad 63$ cf 18 | $18 \quad 32 \quad 63 \mathrm{cf}$ |  | $86 \quad 2618 \quad 76$ |
| ff 3c e5 b0 |  |  |  |  |
| $3 \mathrm{f} 53 \mathrm{6a} 7 \mathrm{c}$ |  |  |  |  |
| 08 e1 ab b7 |  |  |  |  |
| $9 \mathrm{e} \quad 14 \mathrm{7b}$ b9 |  |  |  |  |


[^0]:    * Presenter: Email: salimmw@eng.ukm.my
    ${ }^{\dagger}$ Presenter: Email: nzainal@hotmail.com

